

Math 582

Introduction to Set Theory

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Logic, the language of mathematics

- Ordinary mathematical exposition uses an informal mixture of English and logical notation.
- There is nothing “deep” about such notation: it is just a convenient abbreviation which sometimes increases clarity (and sometimes does not.)
- Our exposition of logical syntax will be **semi-formal** – enough detail to give a reasonably precise account of what **properties** are expressible in the language of set theory.
- A formal development of logical syntax is carried-out in a Logic course.

LAST, the language of set theory

- LAST (LAnguage of Set Theory) is a language suitable for describing mathematical collections (i.e. sets).
- The language will have a precisely determined set of symbols (words) and a rigid syntax (grammar). It is an example of a **formal language** (although my description will be informal).
- The role for this language will be to give a precise and rigorously defined concept of a **definite property**.
- LAST is very simple, but still sufficiently powerful enough to allow that **any set** we may require in mathematics is describable in LAST (i.e. the set consisting of all objects satisfying some definite property describable in LAST).

Variables

☞ We will have symbols for **arbitrary sets**, called variables.

$$V_0, V_1, V_2, \dots$$

☞ A **variable** in mathematics is used like a **pronoun** is used in English. It has no fixed reference (like a **name**), so its meaning “floats” from context-to-context, like a pronoun.

☞ We will have no (officially sanctioned) **names** of sets in LAST.

Relation symbols

☞ We need to be able to make simple assertions about sets. We introduce the following relation symbols:

- Membership symbol: \in ,
- Equality symbol: $=$.

☞ Equality has its intuitive meaning.

☞ However, we will provide **axioms** which will tell you all you need to know about sets and set membership in order to reason about them.

Boolean connective symbols

☞ We need to be able to combine any finite number of simple assertions to produce one big assertion. We use the **Boolean connective symbols**:

$$\vee, \wedge, \neg, \rightarrow, \leftrightarrow$$

☞ The intended meaning and use of these symbols mirrors corresponding English expressions (at least as mathematicians use them):

- \vee means "or",
- \wedge means "and"
- \neg means "not"
- \rightarrow means "implies" ("if-then")
- \leftrightarrow means "iff" ("if and only if")

Quantifier symbols

☞ We will also require two **quantifier symbols**:

- \forall means “for all”,
- \exists means “there exists”

☞ The role of the quantifiers is to provide a **context** to anchor the meaning of the free floating variables.

Punctuation symbols

☞ There will be two punctuation symbols to ensure that the meaning of expressions are unambiguous:

- $(,)$ (parentheses)

Syntax

☞ An **expression** is simply a string of symbols of LAST. Some expressions are more significant, called **formulas**. Here is their definition.

(a) Any expression of the following forms are formulas of LAST.

$$v_n = v_m \quad v_n \in v_m.$$

(b) If φ and ψ are formulas of LAST, then so are

$$(\varphi \wedge \psi) \quad (\varphi \vee \psi) \quad (\varphi \rightarrow \psi) \quad (\varphi \leftrightarrow \psi) \quad (\neg\varphi).$$

(c) If φ is a formula of LAST, then so are

$$\forall v_n \varphi \quad \exists v_n \varphi.$$

(d) Nothing else is a formula of LAST

Examples of Formulas

☞ We establish that an expression is a formula by showing how it can be **constructed** using the rules.

Example. $(v_1 \in v_2 \vee v_2 \in v_1)$ is a formula.

- 1 $v_1 \in v_2$ and $v_2 \in v_1$ are formulas by (a),
- 2 $(v_1 \in v_2 \vee v_2 \in v_1)$ is a formula by (b).

Example. $\forall v_0 \forall v_1 v_0 = v_3$ is a formula.

- 1 $v_0 = v_3$ is a formula by (a),
- 2 $\forall v_1 v_0 = v_3$ is a formula by (c),
- 3 $\forall v_0 \forall v_1 v_0 = v_3$ is a formula by (c).

Unique readability

☞ Our rules of grammar for constructing formulas of LAST is **unambiguous**, in the following sense: there is only one way for applying the rules to construct a given formula. (This is called **unique readability**.)

☞ A formula ψ is called a **subformula** of ϕ if it occurs in ϕ . (So, ϕ is constructed by applying a rule using ψ .)

Examples of Subformulas

Example. The subformulas of $(v_1 \in v_2 \vee v_2 \in v_1)$ are

$$v_1 \in v_2, v_2 \in v_1, (v_1 \in v_2 \vee v_2 \in v_1).$$

Example. The subformulas of $\forall v_0 \forall v_1 v_0 = v_3$ are

$$v_0 = v_3, \forall v_1 v_0 = v_3, \forall v_0 \forall v_1 v_0 = v_3.$$

Free vs. Bound variables

☞ Variables can be either **free** (floating) or **bound**.

- An occurrence of a variable v_n in a formula φ is **free** if v_n does not occur in a subformula of φ of the form $\forall v_n \psi$ or $\exists v_n \psi$. (The book uses the term **parameter** to mean **free variable**.)
- Otherwise, v_n is said to be **bound**

☞ A formula of LAST with no free variables is a **sentence**.

The sentences of LAST have a fully determinate meaning, so can be used to make assertions.

☞ A formula with free variables will have an indeterminate meaning, since its free variables will have no fixed determinate meaning.

Examples of free and bound variables

Example. $\forall v_1 (v_1 \in v_2 \rightarrow \exists v_3 v_3 \in v_2)$.

- free: v_2 (both occurrences),
- bound: v_1, v_3 .

Example. $\forall v_2 (v_1 \in v_2 \rightarrow \exists v_1 v_1 \in v_2)$.

- free: v_1 (first occurrence only),
- bound: v_1 (second occurrence only), v_2 (both occurrences).

Example. Neither of the previous examples were sentences.

$\exists v_4 \forall v_3 (v_4 \in v_3 \rightarrow v_3 \in v_4)$ is a sentence.

Definite Properties

☞ We will write $\varphi(v_0, \dots, v_n)$ to indicate that φ is a formula whose free variables, if any, are among the list v_0, \dots, v_n .

☞ A **definite property** in LAST is any formula of the form $\varphi(v_0)$.
A **definite relation** in LAST is any formula of the form $\varphi(v_0, \dots, v_n)$.

Examples.

- $v_0 \in v_1$ (the membership relation),
- $\exists v_1 v_1 \in v_0$ (v_0 is a nonempty set),
- $v_0 \notin v_0$ (v_0 is not a member of itself),
- $\exists v_1 (v_1 \in v_0 \wedge \forall v_2 (v_2 \in v_0 \rightarrow v_2 = v_1))$ (v_0 is a singleton set).

Conventions

- We will use a wide variety of typographic symbols to stand for variables in our language:
 $x, y, z, X, Y, Z, a, b, c, A, B, C, \alpha, \beta, \gamma, \kappa, \lambda, \mathcal{F}$. The reason is that it will make it easier to read statements if you know that α is always an ordinal number, κ a cardinal number, \mathcal{F} is a family of sets, etc.
- Certain variable expressions, “ φ ”, “ ψ ”, are not part of the language, but are used to **talk about** expressions of the language. They are part of the **metalanguage** – the language used to talk about our mathematical language.
- The text uses capital **boldface**: **P**, **A**, etc. to refer to syntactic expressions (sentences and formulas.) They write **P**(x) where I use $\varphi(x)$.
- We will use brackets: [,] as well as parentheses if this improves readability.