

A CHARACTERIZATION OF THE low_n DEGREES

KENNETH HARRIS

Department of Computer Science
University of Chicago

<http://people.cs.uchicago.edu/~kaharris>

kaharris@uchicago.edu

Dominate and Escape

Let $f, g : \mathbb{N} \rightarrow \mathbb{N}$.

- f *dominates* g if

$$(\forall^\infty x) [f(x) > g(x)]$$

f is a *dominant function* if f dominates every computable function.

- g *escapes* (domination from) f if

$$(\exists^\infty x) [f(x) \leq g(x)]$$

f has an *escape function* if there is a computable g which escapes domination by f .

Martin's Characterization

Theorem (*Martin*, 1966) Let \mathbf{a} be a Turing degree.

- \mathbf{a} is high ($\mathbf{a}' \geq \mathbf{0}''$) iff there is an \mathbf{a} -computable dominant function:

$$(\exists f \leq \mathbf{a})(\forall g \leq \mathbf{0}) [f \text{ dominates } g]$$

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- \mathbf{a} is non-high ($\mathbf{a}' < \mathbf{0}''$) iff every \mathbf{a} -computable function has an escape function:

$$(\forall f \leq \mathbf{a})(\exists g \leq \mathbf{0}) [g \text{ escapes } f]$$

Uniform Escape Property

Question: For what non-high degrees can escape functions be *effectively* produced?

Definition: A degree \mathbf{a} has the *Uniform Escape Property* (UEP), or (1-UEP), when for any set $A \in \mathbf{a}$:

There is a partial computable $\lambda ex.h_e(x)$ such that whenever Φ_e^A is total, then

h_e total and escapes Φ_e^A

Recall, h_e escapes Φ_e^A if

$$(\exists^\infty x) [\Phi_e^A(x) \leq h_e(x)]$$

UEP Equivalent to low_1

Theorem: For all degrees \mathbf{a} TFAE

- (A) \mathbf{a} is low_1 ($\mathbf{a}' \leq \mathbf{0}'$).
- (B) \mathbf{a} has the Uniform Escape Property.

low_n Degrees and Escape Functions

There is a hierarchy of properties characterized by progressively less effective procedures, *n-Uniform Escape Property (n-UEP)*, starting with (1-UEP)=(UEP), such that

Theorem: For all degrees \mathbf{a} and all $n \geq 1$ TFAE

- (A) \mathbf{a} is low_n ($\mathbf{a}^{(n)} \leq \mathbf{0}^{(n)}$).
- (B) \mathbf{a} has (*n-UEP*).

Quantifiers on Steroids

$(\forall^\infty x)$: *For almost every x .*

Reduces to $\exists\forall$ and behaves like \forall .

$(\exists^\infty x)$: *There exists infinitely many x .*

Reduces to $\forall\exists$ and behaves like \exists .

Fundamental Relations

$$\exists^\infty P \iff \neg\forall^\infty\neg P$$

$$\forall \implies \forall^\infty \implies \exists^\infty \implies \exists$$

Theorem (Strong Normal Form):

The arithmetic hierarchy is characterized by alternations of the two strongest quantifiers, \forall and \forall^∞ .

low_1 DEGREES AND
THE 1-UNIFORM ESCAPE PROPERTY

low_1 implies Uniform Escape Property

Theorem: All low_1 sets A have (1-UEP):

There is a partial computable function $\lambda x.h_e(x)$
such that whenever Φ_e^A is total, then

h_e total and escapes Φ_e^A

The Key Idea

Let A be low₁, so $\Pi_2^A \subset \Pi_2$.

Want: Computable g such that for each total Φ_e^A , $W_{g(e)}$ satisfies

(**escape**) $(\exists^\infty x)(\exists s) [\Phi_{e,s}^A(x) \downarrow \leq s \ \& \ x \notin W_{g(e),s}]$

(**total**) $W_{g(e)} = \omega$

Define:

$$h_e(x) = (\mu s) [x \in W_{g(e),s}]$$

Problem: How to match (**escape**) with (**total**)?

Strong Normal Form: Π_2, Σ_2

Normal Form: For $V \in \Pi_2$, there is some v (the Π_2 index for V) with

$$V(e) \iff (\forall y) [\langle e, y \rangle \in W_v]$$

Strong Normal Form (SNF): There is a computable g , such that for any $V \in \Pi_2$ with index e

$$\begin{aligned} V(e) &\implies [W_{g(v,e)} = \omega] \\ \neg V(e) &\implies [W_{g(v,e)} \text{ finite}] \end{aligned}$$

Implementation of Key Idea

Let A be low₁ (thus $\Pi_2^A \subset \Pi_2$).

Define Π_2^A predicate (**escape**)

$$(\exists^\infty x)(\exists s) [\Phi_{e,s}^A(x) \downarrow \leq s \ \& \ x \notin W_{g(v,e),s}]$$

where g is the computable function given by (SNF) from a Π_2 index v for (**escape**):

$$\begin{aligned} (\mathbf{escape}) &\implies [W_{g(v,e)} = \omega] \\ \neg(\mathbf{escape}) &\implies [W_{g(v,e)} \text{ finite}] \end{aligned}$$

then define

$$h_e(x) = (\mu s) [x \in W_{g(v,e),s}]$$

low_2 DEGREES AND
THE 2-UNIFORM ESCAPE PROPERTY

2-Uniform Escape Property: First Change

Definition: A set A is low_2 if $A'' \leq \mathbf{0}''$.

With low_2 we add one jump class and one layer of quantifier complexity.

Our *first change* in defining (2-UEP):

There are uniformly enumerable (u.e.) families of partial computable functions $\lambda e. \{h_{e,y}\}_{y \in \omega}$ such that whenever Φ_e^A is total, then

$$(\forall^\infty y) \left[h_{e,y} \text{ total and escapes } \Phi_e^A \right]$$

2-Uniform Escape Property

Definition: A degree \mathbf{a} has the *2-Uniform Escape Property* (2-UEP), when for any set $A \in \mathbf{a}$:

There are uniformly enumerable (u.e.) families of partial computable functions $\lambda e. \{h_{e,y}\}_{y \in \omega}$ such that for any u.e. family of functions $\{\Phi_{e,y}^A\}_{y \in \omega}$ satisfying

$$(\forall^\infty y) [\Phi_{e,y}^A \text{ total}]$$

then

$$(\forall^\infty y) \left[h_{e,y} \text{ total and escapes } \Phi_{e,y}^A \right]$$

low_2 Equivalent to 2-UEP

For all degrees \mathbf{a} TFAE

- (A) \mathbf{a} is low_2 ($\mathbf{a}'' \leq \mathbf{0}''$).
- (B) \mathbf{a} has the 2-Uniform Escape Property.

Strong Normal Form: Σ_3, Π_3

If A is low₂ then $\Sigma_3^A \subseteq \Sigma_3$.

Strategy of Proof: Pump (**escape**) property (Π_2^A) with strong quantifiers to Σ_3^A and exploit weakness of A .

$$(\forall^\infty y)(\exists^\infty x)(\exists s) [\Phi_{e,y,s}^A(x) \downarrow \leq s \ \& \ x \notin W_{g(v,e,y),s}]$$

Strong Normal Form (SNF): There is a computable g , such that for any $V \in \Sigma_3$ with index e

$$\begin{aligned} (\text{escape}) &\implies (\forall^\infty y) [W_{g(v,e,y)} = \omega] \\ \neg(\text{escape}) &\implies (\forall y) [W_{g(v,e,y)} \text{ finite}] \end{aligned}$$

low_3 DEGREES AND BEYOND

3-Uniform Escape Property

A degree \mathbf{a} is low₃ if $\mathbf{a}''' = \mathbf{0}'''$.

Definition: A degree \mathbf{a} has the *3-Uniform Escape Property* (3-UEP) when for any set $A \in \mathbf{a}$:

There are uniformly enumerable (u.e.) families of partial computable functions $\lambda e. \{h_{e,y_1,y_2}\}_{y_1,y_2 \in \omega}$ such that for any u.e. family of functions

$\{\Phi_{e,y_1,y_2}^A\}_{y_1,y_2 \in \omega}$ satisfying

$$(\exists^\infty y_2)(\forall^\infty y_1) [\Phi_{e,y_1,y_2}^A \text{ total}]$$

then

$$(\exists^\infty y_2)(\forall^\infty y_1) \left[h_{e,y_1,y_2} \text{ total and escapes } \Phi_{e,y_1,y_2}^A \right]$$

low_3 Equivalent to 3-UEP

Theorem: For all degrees \mathbf{a} TFAE

- (A) \mathbf{a} is low_3 ($\mathbf{a}''' \leq \mathbf{0}'''$).
- (B) \mathbf{a} has the 3-Uniform Escape Property.

Strong Normal Form: Π_4, Σ_4

If A is low₃ then $\Pi_4^A \subseteq \Pi_4$.

Strategy of Proof: Pump (**escape**) property (Π_2^A) with strong quantifiers to Π_4^A and exploit weakness of A .

$$(\exists^\infty y_2)(\forall^\infty y_1)(\exists^\infty x)(\exists s) \left[\Phi_{e,y_1,y_2,s}^A(x) \downarrow \leq s \right. \\ \left. \& x \notin W_{g(v,e,y_1,y_2),s} \right]$$

Strong Normal Form (SNF): There is a computable g , such that for any $V \in \Pi_4$ with index e

$$\begin{aligned} (\mathbf{escape}) &\implies (\forall y_2)(\forall^\infty y_1) \left[W_{g(v,e,y_1,y_2)} = \omega \right] \\ \neg(\mathbf{escape}) &\implies (\forall^\infty y_2)(\forall y_1) \left[W_{g(v,e,y_1,y_2)} \text{ finite} \right] \end{aligned}$$

n-Uniform Escape Property

A degree \mathbf{a} is low_n if $\mathbf{a}^{(n)} = \mathbf{0}^{(n)}$.

Definition: A degree \mathbf{a} has the *n-Uniform Escape Property* (*n-UEP*) when for any set $A \in \mathbf{a}$:

There are uniformly enumerable (u.e.) families of partial computable functions $\lambda e. \{h_{e,\bar{y}}\}_{\bar{y} \in \omega}$ such that for any u.e. family of functions $\{\Phi_{e,\bar{y}}^A\}_{\bar{y} \in \omega}$ satisfying

$$(Q_1 y_{n-1})(Q_2 y_{n-2}) \dots [\Phi_{e,\bar{y}}^A \text{ total}]$$

then

$$(Q_1 y_{n-1})(Q_2 y_{n-2}) \dots \left[h_{e,\bar{y}} \text{ total and escapes } \Phi_{e,\bar{y}}^A \right]$$

where $Q_1, Q_2 \in \{\exists^\infty, \forall^\infty\}$ by

- For *odd* n : alternate $\exists^\infty \forall^\infty$
- For *even* n : alternate $\forall^\infty \exists^\infty$

low_n Equivalent to n -UEP

Theorem: For all degrees \mathbf{a} TFAE

- (A) \mathbf{a} is low_n ($\mathbf{a}^{(n)} \leq \mathbf{0}^{(n)}$).
- (B) \mathbf{a} has the n -Uniform Escape Property.

Strong Normal Form Theorem

Strong Normal Form Theorem (SNF) (with $n \geq 1$)

All arithmetic formulas equivalent to formulas using only the *beefiest* quantifiers $\{\forall, \forall^\infty\}$:

For any $V \in \Sigma_{2n+1}$ with index v there is a computable g , such that

$$\begin{aligned} V(e) &\implies (\forall^\infty y_{2n-1})(\forall y_{2n-2}) \dots [W_{g(v,e,\bar{y})} = \omega] \\ \neg V(e) &\implies (\forall y_{2n-1})(\forall^\infty y_{2n-2}) \dots [W_{g(v,e,\bar{y})} \text{ finite}] \end{aligned}$$

For any $U \in \Pi_{2n}$ with index u there is a computable g , such that

$$\begin{aligned} U(e) &\implies (\forall y_{2n-2})(\forall^\infty y_{2n-3}) \dots [W_{g(u,e,\bar{y})} = \omega] \\ \neg U(e) &\implies (\forall^\infty y_{2n-3})(\forall y_{2n-2}) \dots [W_{g(u,e,\bar{y})} \text{ finite}] \end{aligned}$$

APPLICATION OF ESCAPE FUNCTIONS

Bounding Saturated Models

Theorem: There is a complete decidable theory \mathcal{T} whose types are all computable, which has *no* saturated model of low_n c.e. degree for any n .

Bibliography

My work:

Kenneth Harris, "A Characterization of the low_n Degrees using Escape Functions, preprint at people.cs.uchicago.edu/~kaharris/papers/lown.pdf

Kenneth Harris, "On Bounding Saturated Models", preprint at people.cs.uchicago.edu/~kaharris/papers/sat.pdf